

# Enumeration Algorithms for *Restricted and Unrestricted* Compositions and Words

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# Enumeration?

Two very different disciplines

- Counting
- Generation

Two different disciplines in computation but a generation solution is more powerful.

*Typical interests:*

- **Combinatorial Objects**
- **Discrete Solution Spaces**
- **Sequencing**
- **Many more...**

# Enumeration?

Primary Goals of an enumeration algorithm:

- Given properties of a finite set, exhaust all the members of the set which fit the definition.
- Avoiding extra structure whenever possible.
- Understanding more about the structure of a construction.
- Can solve counting problems.
- Eg) Combinatorial Objects, Word Problems...

**WARNING: These algorithms are highly output sensitive.**

# Combinatorial Composition

Definition: A finite set  $C_{s,n}$  which contains all non-negative integer sequences  $\lambda = \lambda_1 \lambda_2 \dots \lambda_n$  such that  $\sum (\lambda_i) = s$ ,  $1 \leq i \leq n$ . This is called an *unrestricted composition*.

e.g)  $C_{3,3} = \{(0,0,3), (0,3,0), (3,0,0), (2,1,0), (1,2,0), (0,2,1), (0,1,2), (1,0,2), (2,0,1), (1,1,1)\}$

# Combinatorial Composition

There is also a more general version:

A finite set  $C_{s,n}^R$  which contains all non-negative integer sequences  $\lambda = \lambda_1 \lambda_2 \dots \lambda_n$  such that  $\sum (\lambda_i) = s$ ,  $1 \leq i \leq n$ , where  $\lambda_i \in R$ .

This is called an *restricted composition*. There are several types of restricted compositions...

e.g)  $C_{3,3}^R = \{(1,1,1)\}$ , where  $R = \{1\}$

# History

Compositions play a vital role in the foundations of Theoretical Computer Science and also are of interest in Combinatorics.

Generation of Compositions are a specific subtype of Word Problems.

Word Problem – Given a semi-Thue System,

- rewrite a string given a set of 'rewriting' axioms, to produce an outputting string in the string rewriting system in a language.
- General Case: Problem is *undecidable*.

# History

## Early Research:

Percy Alexander MacMahon (1854-1929)



*“Compositions are merely partitions in which the order of occurrence of the parts are important”*

- Many of his conventions are still used today.
- Earliest results.
- *“Assemblage of objects”*

Credit: “Combinatorics of Compositions and Words”, Heubach,Mansour

# History

## Early Research:

Axel Thue (1863-1922)



*“Zeichenreihen”*

- First to systematically study words and compositions.
- Extended to the infinite cases. The general word problem.
- Thue and semi-Thue systems.

Credit: “Combinatorics of Compositions and Words”, Heubach, Mansour

*“Combinatorics on Words”*

# Problems

## Problem 1:

Given positive integers  $n, s$  ( $n = 0$  iff  $s = 0$ ),  
enumerate the unrestricted composition  
 $C_{s,n}$ .

- Some Results:
- $C(s,n) \stackrel{s+n-1}{=} \text{CHOOSE}(s+n-1, n-1)$
- Lower bound space-complexity:  $\Omega(C(s,n))$ .
  - Output sensitive
  - Tight bound (it's the output itself).

# Problems

## Output Sensitivity Extremes:

Keep this in mind! This is your solution space size that you output...

- $C(8,7) = 3003.$
- $C(9,9) = 24310.$
- $C(200,100) = 1.38608382108618824826 \times 10^{81}$

# Problems

## Some Solutions to Problem 1:

Often solutions try to take one element and form the next from the previous to avoid consuming extra space.

### Randomized Generation

- Idea: Keep generating unique sequences of unique sum of  $s$  from a uniform distribution on  $[0,s]$  and test if it's currently in the set.
- Hard to analyze, but often useful in parallel implementations.

# Problems

## Some Solutions to Problem 1:

### K-bit Reflected Gray Code Generation:

- Idea: Generate all gray codes in lexicographical order of length  $n$  integers in binary, then on the last iteration, remove all elements which don't sum to  $s$  if the last symbol is added.
- Constructs elements from scratch but the trade off could be dependent on implementation. (moves to modern results)
- Bitner-Ehrlich-Reingold (BER) method.

# Problems

## Some Solutions to Problem 1:

### Brute Force:

- Idea: Similar to the previous technique except working directly with integers.
- Based on the definition, this technique can always take advantage of the fact we have a continuous interval of non-negative integers.
- Simpler to implement but space can grow substantially more than the previous. Very small instances can be efficient.

# Problems

## Upper Bounds/Lower Bounds:

This problem has been attacked numerous times by many researchers since the 70's by the likes of Knuth, Lothaire, and many others.

Optimal solution: Combinatorial Gray Codes.

For simplicity: let  $k$  be the output size of a given composition.

# Problems

## Vital Result:

Goal: Compute gray codes quickly, and perform the 2<sup>nd</sup> technique (with modification).

- O(1) – Generating gray codes in O(1) worst-case time per word [Walsh,2003]
- Map this problem to our unrestricted composition problem
- Upper Bound –  $\Omega(k)$  [evident from above]
- Optimal Time.

# Problems

## Problem 2 (Definition):

A finite set  $C^{L,U}_{s,n}$  which contains all integer sequences  $\lambda = \lambda_1 \lambda_2 \dots \lambda_n$  such that  $\sum (\lambda_i) = s$ ,  $1 \leq i \leq n$ , where  $L \leq \lambda_i \leq U$ .

This is called an *bounded composition*.

- *Specific type of restricted composition.*
- Inherits problems from problem 1.
- New problems: unused space and decision complexity increases.

# Problems

## Problem 2:

Given positive integers  $n, s$  ( $n = 0$  iff  $s = 0$ ), and non-negative integers  $L, U$ , enumerate the bounded composition  $C_{s,n}^{L,U}$ .

- *Why would we want this?*
  - *Many of the same reasons, but some of the more modern research was for statistical analysis in computation.*
- Applying the last problems solutions can be tedious but are doable but can lead to space-issues in contrast to solution size.

# Problems

## Problem 2:

Using the same techniques as Problem 1 would cause a rise in time-complexity using the optimal solution.

- When  $L=0$  and  $U=s$ , we get the unrestricted composition.
- Same lower bound holds since this is a sub-instance of Problem 1.

# Problems

## Some Solutions to Problem 2:

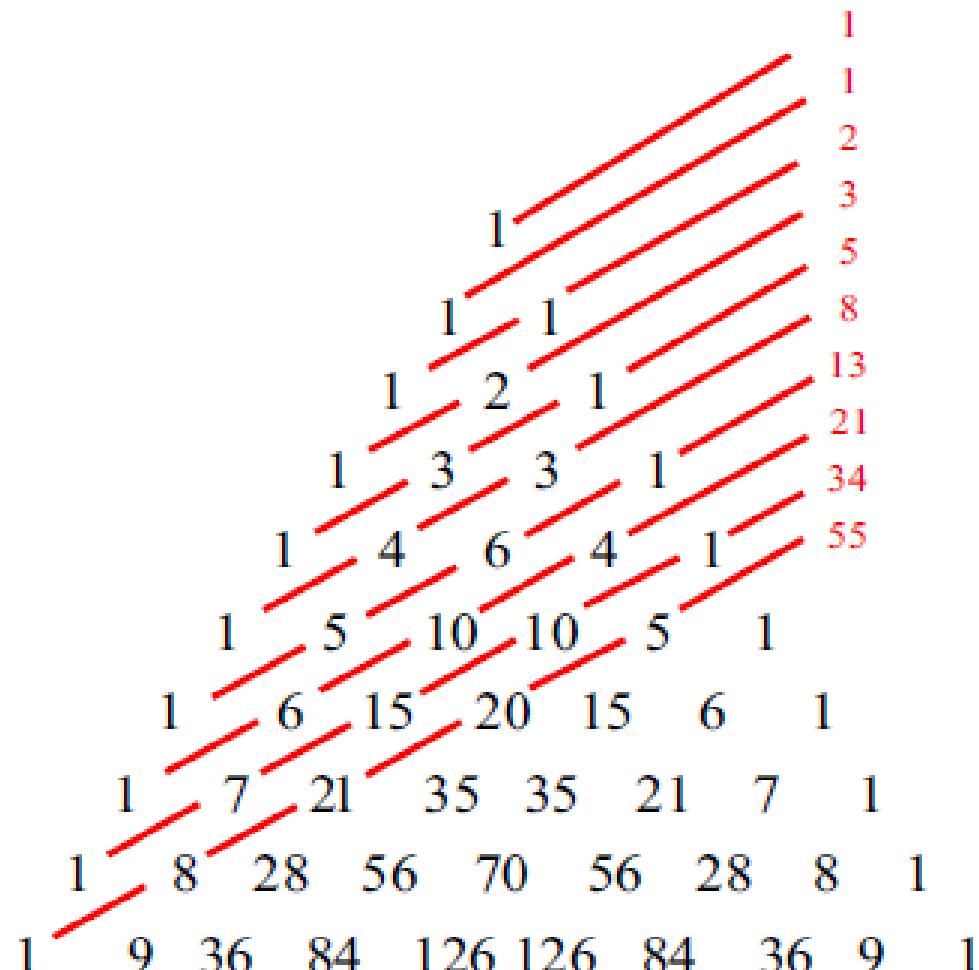
Another technique:

Idea: Based on Fibonnaci Numbers and  
Pascal's Triangle for counting.[Kimberling.  
2002]

- Inside pascal's triangle there exists a path from an interior node to the root that will diagonally reach the root of the triangle. Thus when reading off the coefficients of the triangle you obtain a composition sequence of elements for each off diagonal.
- Picture... next here

# Problems

Some Solutions to Problem 2:



credit: JD Opdyke, 2010

# Problems

## Some Solutions to Problem 2:

$$\begin{array}{ccccc} \binom{0}{0} & & & & \\ \binom{1}{0} & \binom{1}{1} & & & \\ \binom{2}{0} & \binom{2}{1} & \binom{2}{2} & & \\ \binom{3}{0} & \binom{3}{1} & \binom{3}{2} & \binom{3}{3} & \\ \binom{4}{0} & \binom{4}{1} & \binom{4}{2} & \binom{4}{3} & \binom{4}{4} \\ \binom{5}{0} & \binom{5}{1} & \binom{5}{2} & \binom{5}{3} & \binom{5}{4} & \binom{5}{5} \\ \binom{6}{0} & \binom{6}{1} & \binom{6}{2} & \binom{6}{3} & \binom{6}{4} & \binom{6}{5} & \binom{6}{6} \\ \binom{7}{0} & \binom{7}{1} & \binom{7}{2} & \binom{7}{3} & \binom{7}{4} & \binom{7}{5} & \binom{7}{6} & \binom{7}{7} \\ \binom{8}{0} & \binom{8}{1} & \binom{8}{2} & \binom{8}{3} & \binom{8}{4} & \binom{8}{5} & \binom{8}{6} & \binom{8}{7} & \binom{8}{8} \\ \binom{9}{0} & \binom{9}{1} & \binom{9}{2} & \binom{9}{3} & \binom{9}{4} & \binom{9}{5} & \binom{9}{6} & \binom{9}{7} & \binom{9}{8} & \binom{9}{9} \end{array}$$

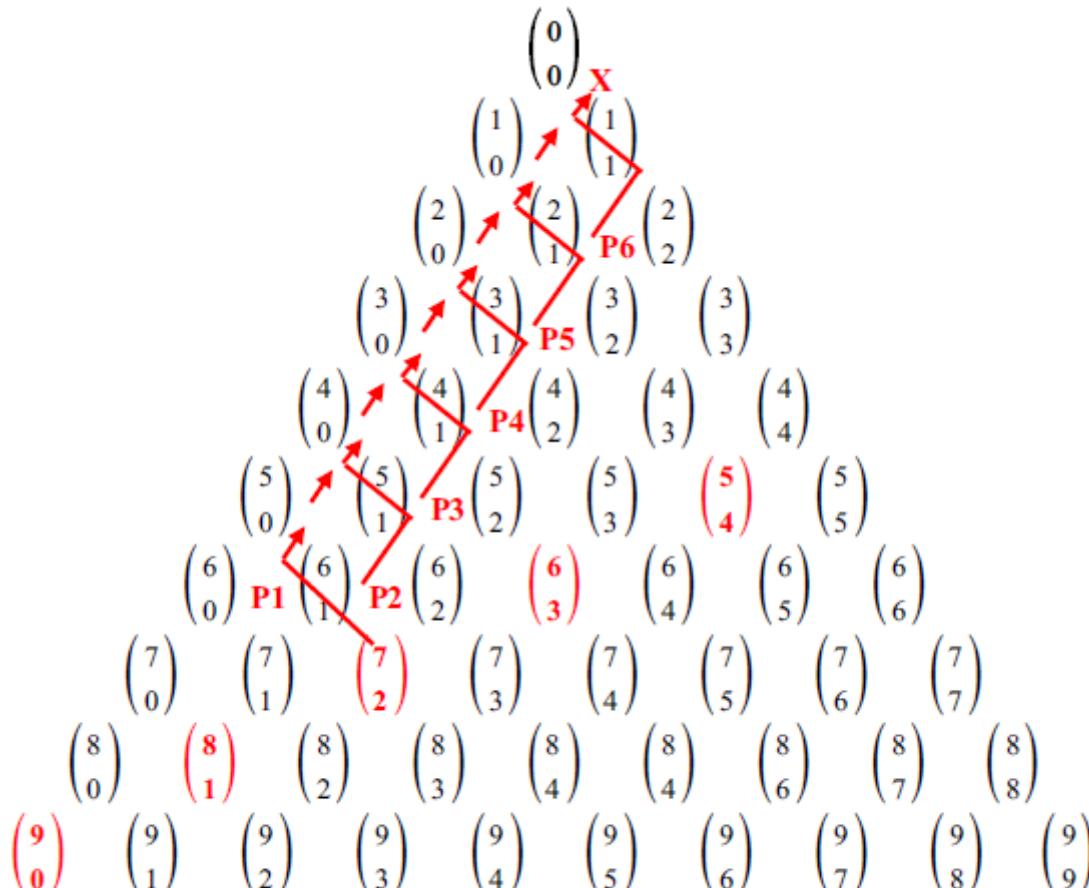
$$n = 11 = (9 + 2) = (\text{row\#} + 2)$$

$$\max(k) = \lfloor n/a \rfloor = \lfloor 11/2 \rfloor = 5 \text{ off-diagonals}$$

credit: JD Opdyke, 2010

# Problems

## Some Solutions to Problem 2:



Composition paths traced by RICs\_Base( $n = 11, k = 3$ )

credit: JD Opdyke, 2010

# Problems

## Some Solutions to Problem 2:

This solution was developed by J.D. Opdyke in 2010 from Kimberling's result.

- time-complexity:  $\sim O(k \times \text{RICs})$
- Hard to bound the algorithm's run-time.
- Flaws: Recursive, and has loops.
- Is the best result in Bounded Compositions.

$$O\left(\sum_{k=\min k}^{\max k} \left( k \cdot \sum_{i=(n-b)}^{(n-a)} c(i, k-1, a, b) \right)\right)$$

# Future Problems

## Problem 3:

Given positive integers  $n, s$  ( $n = 0$  iff  $s = 0$ ), and set of non-negative integers  $R$ , enumerate the restricted composition  $C_{s,n}^R$ .

- Very tough problem. General 'first order'.
- Has reduction to Positive Integer Subset Sum
- Problem 1 and Problem 2's solutions won't work on this without a heavy cost.  
Discontinuity of values.
- Work on counting has been done for this but, not generating.

# Future Research

## Research Interests:

We had mentioned Problem 3 but what other interests hold in this discourse with Compositions still?

- Pattern avoidance work, w-avoidance
- Relationships between Integer Partitions and Integer Compositions
- '2<sup>nd</sup> order' problems with compositions.

# Questions?

Thank you very much for your time!

Questions?

*Have a beautiful day!*